

# Hybrid Methods for global optimization of large-scale polynomials

J. Aspman<sup>1</sup> G. Bareilles<sup>1</sup> V. Kungurtsev<sup>1</sup> J. Mareček<sup>1</sup> M. Takáč<sup>2</sup>

<sup>1</sup>Czech Technical University <sup>2</sup>Mohamed bin Zayed Univ. of A.I.

## Problem & contributions

$$\begin{aligned} f^* = \min_x f(x) \\ \text{s.t. } g_i(x) \geq 0, \quad i \in \mathcal{I} \\ g_j(x) = 0, \quad j \in \mathcal{E} \end{aligned} \quad (\text{POP})$$

where  $f, g_i, g_j : \mathbb{R}^n \rightarrow \mathbb{R}$  are polynomials.

▷ Problem

Find a *global* minimizer  $x^*$  of large-scale (POP)

▷ Contribution

- bring together tools from different math disciplines
- a *hybrid* algorithm
  - a 1st order method on a convex relaxation ( $P_r$ )
  - a theoretically grounded switch to Newton's method on (POP)

## Difficulties & approach

The global minimizer of large (POP) is provided by very large SDP relaxations ( $P_r$ ).

▷ Interior Point Methods: go-to method for solving ( $P_r$ )

- ✓ fast convergence
- ✗ not applicable (too high memory cost)

▷ First-order methods on ( $P_r$ ) – fig. 2

- ✓ low per-iteration cost
- ≈ slow convergence: provide rough solutions of ( $P_r$ )

▷ Newton on smooth polynomial equations – fig. 1

- ≈ poor convergence away from solution
- ✓ superfast local convergence near solutions

▷ Active constraint identification – fig. 3

- ✓ turns (POP) into smooth poly. eq. near  $x^*$

## Illustrative problem

$$\begin{aligned} \min_{x \in \mathbb{R}^2} f(x) = -2.5x_1^2 + 3x_1x_2 - 2.5x_2^2 - 3x_1 + 5x_2 - 2.5 \\ \text{s.t. } g_1(x) = -0.5x_1^3 + x_2 \geq 0 \\ g_2(x) = -0.05x_1^2 - x_2 + 1.8 \geq 0 \\ g_3(x) = -0.05x_2^2 + x_1 + 0.1x_2 + 0.35 \geq 0 \end{aligned} \quad (1)$$

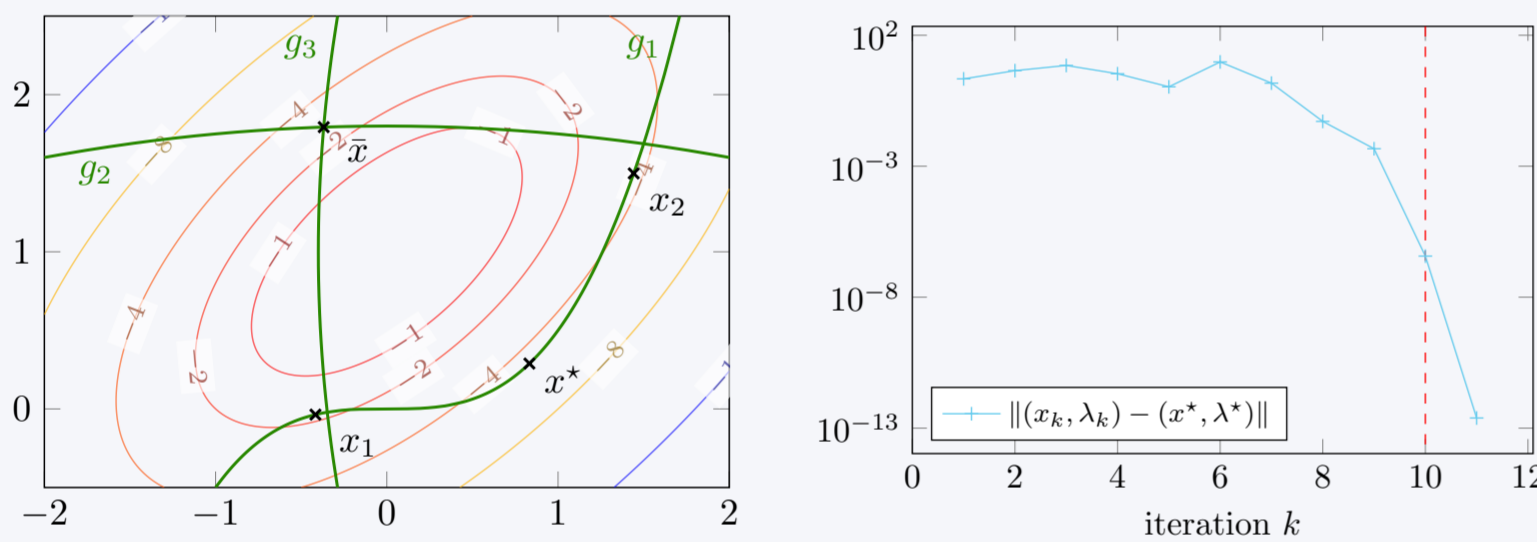


Figure 1. Left pane: problem (1). Right pane: Newton's method: slow away from  $(x^*, \lambda^*)$ , fast near  $(x^*, \lambda^*)$ .

## 1. Convex relaxations

**TLDR:** ( $P_r$ ) is a sequence of increasingly large SDP problems, whose optimal value converge to  $f^*$  from below.

▷ Moment relaxation, order  $r$

$$\begin{aligned} \rho_r = \inf_{y \in \mathbb{R}^{\mathbb{N}_{2r}^n}} L_y(f) \\ \text{s.t. } M_r(y) \geq 0 \\ M_{r-r_j}(g_j y) = 0, \quad j \in \mathcal{E} \\ M_{r-r_j}(g_j y) \geq 0, \quad j \in \mathcal{I} \\ y_0 = 1 \end{aligned} \quad (P_r)$$

where  $r_j = \lceil (\deg g_j)/2 \rceil$ ,  $y$  is a sequence indexed by exponents  $\alpha \in \mathbb{N}_{2r}^n = \{\alpha \in \mathbb{N}^n : \sum \alpha_i \leq 2r\}$ ,  $L_y : f = \sum_{\alpha} f_{\alpha} X^{\alpha} \mapsto \sum_{\alpha} f_{\alpha} y_{\alpha}$  is the **Riesz functional**, and  $M_d(y)$  and  $M_d(g_j y)$  are the **moment** and **localizing matrices**:

$$\begin{aligned} M_d(y)[\alpha, \beta] &= L_y(X^{\alpha} X^{\beta}) = y_{\alpha+\beta} \\ M_d(g_j y)[\alpha, \beta] &= L_y(g_j X^{\alpha} X^{\beta}) = \sum_{\delta} g_{j,\delta} y_{\alpha+\beta+\delta}, \end{aligned}$$

for all  $\alpha, \beta \in \mathbb{N}_{2r}^n$ .

▷ **Property** [3]:  $\rho_r$  converges to  $f^*$  from below.

If ( $P_r$ ) has a unique global minimizer  $x^*$ , and  $y^r$  is a nearly optimal solution of ( $P_r$ ), then

$$\lim_{r \rightarrow \infty} L_{y^r}(X_j) = x_j^*$$

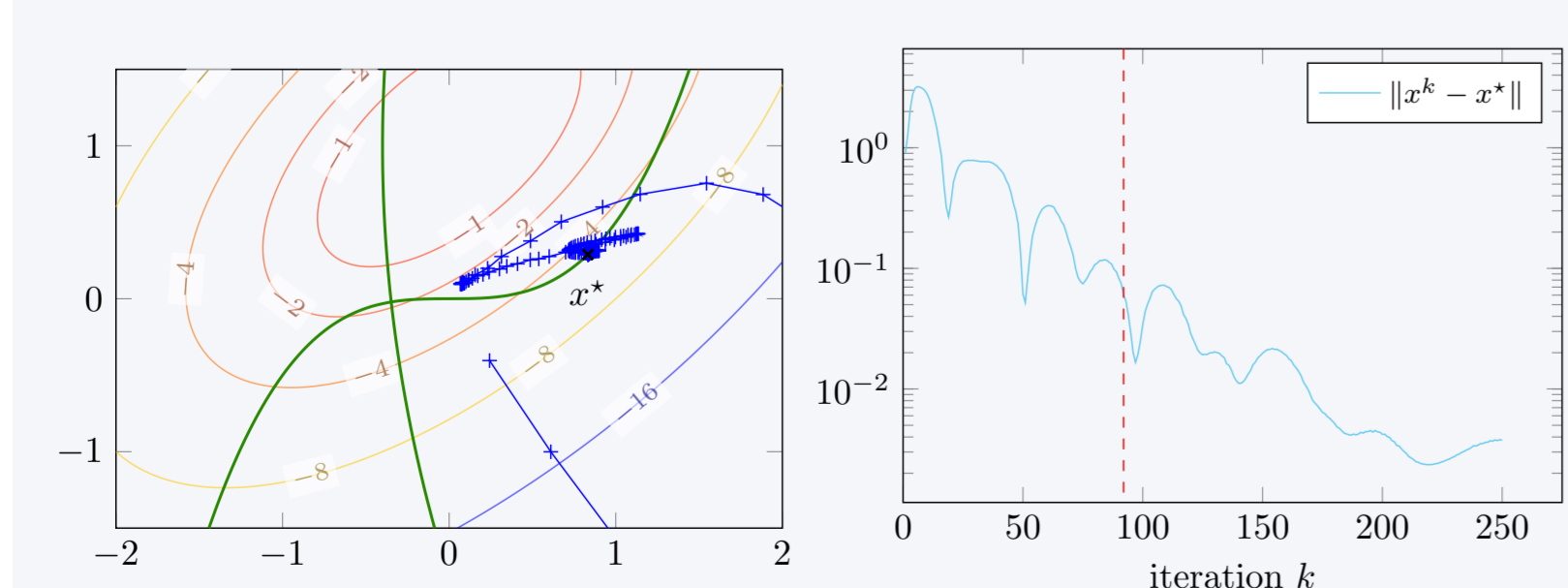


Figure 2. First-order method on the third relaxation ( $P_3$ ) of (1). Left pane: points  $x^k = (L_{y^k}(X_1), L_{y^k}(X_2))_k \in \mathbb{R}^2$ , where  $y^k \in \mathbb{R}^{\mathbb{N}_6^2}$  is the  $k$ -th iterate of the 1st order method on ( $P_3$ ). Right pane: distance between  $x^k$  and the global minimizer  $x^*$ . The dashed red line indicates the first time the correct active set is detected.

## 2. Identifying active constraints

**TLDR:** at  $x^*$ , some constraints are null and others positive.

$$\mathcal{A}^* = \{i \in \mathcal{I} : g_i(x^*) = 0\}$$

can be detected from any point  $x$  near  $x^*$ .

▷ The Lagrangian for (POP) is

$$\mathcal{L}(x, \lambda) = f(x) - \lambda^T g(x),$$

where  $\lambda \in \mathbb{R}^{m_{\mathcal{E}}} \times \mathbb{R}_+^{m_{\mathcal{I}}}$  is the Lagrange multiplier vector.

▷ The first-order “KKT” necessary optimality conditions: there exist  $\lambda^*$  such that

$$\begin{aligned} \nabla_x \mathcal{L}(x^*, \lambda^*) &= 0 \\ g_{\mathcal{E}}(x^*) &= 0 \\ 0 \leq g_{\mathcal{I}}(x^*) \perp \lambda_{\mathcal{I}}^* \geq 0 \end{aligned} \quad (2)$$

where  $a \perp b$  means  $a^T b = 0$ .

▷ Measure of the KKT residual at point  $(x, \lambda)$ :

$$\begin{aligned} r(x, \lambda) = \|\nabla_x \mathcal{L}(x, \lambda)\| + \|g_{\mathcal{E}}(x)\| \\ + \|[-g_{\mathcal{I}}(x)]_+\| + |\lambda_{\mathcal{I}}^T g_{\mathcal{I}}(x)| + \|[-\lambda_{\mathcal{I}}]\| \end{aligned}$$

▷ From point  $x$ , the active set is chosen as

$$\begin{aligned} \omega(x) &= \min_{\lambda \in \mathbb{R}^{m_{\mathcal{E}}} \times \mathbb{R}_+^{m_{\mathcal{I}}}} r(x, \lambda) \\ \mathcal{A}(x) &= \{i : g_i(x) \leq -1/\log(\omega(x))\} \end{aligned}$$

**Theorem** [1, Th. 3.2]: Consider  $x^*$  that satisfies KKT (2), a qualification condition  $\text{MFCQ}$ , and is an isolated critical point. Then, there exists  $\epsilon > 0$  such that

$$\mathcal{A}(x) = \mathcal{A}^* \quad \text{for all } x : \|x - x^*\| < \epsilon.$$

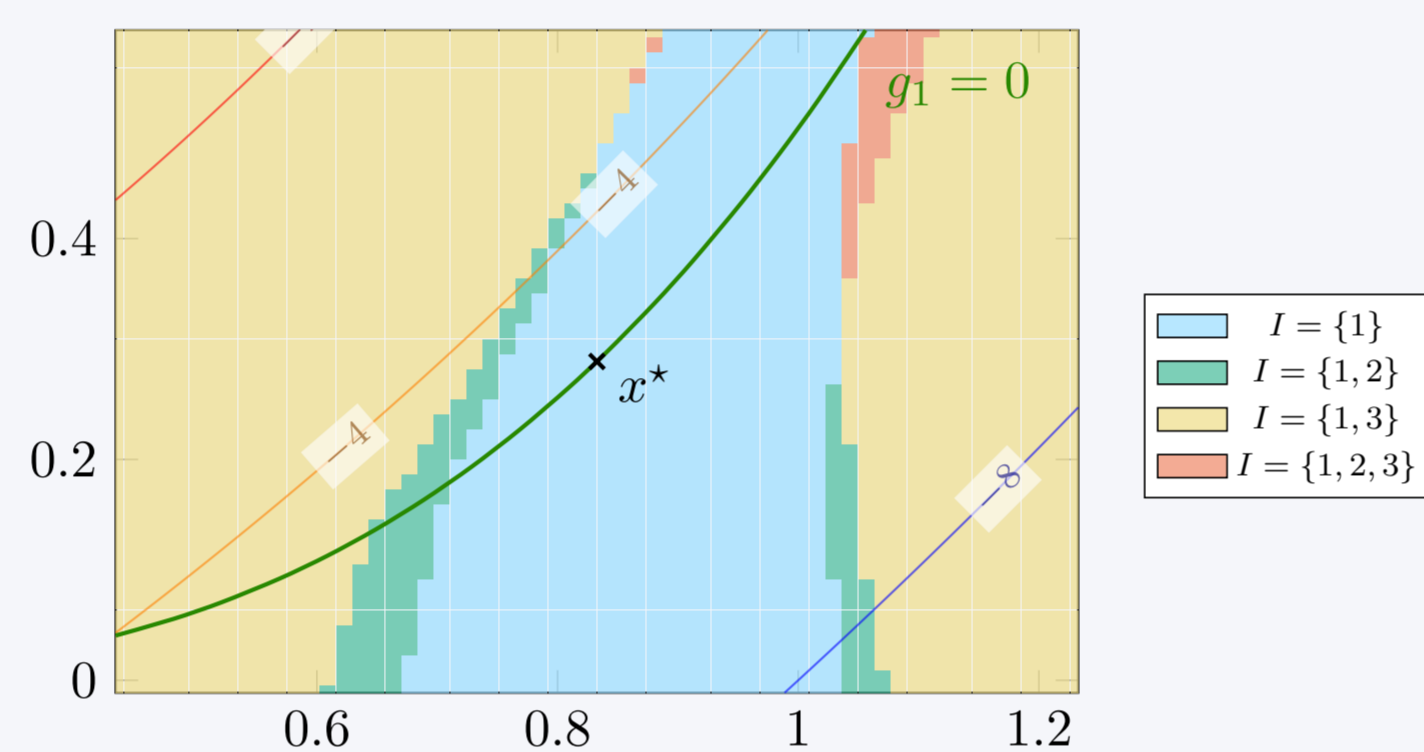


Figure 3.  $\mathcal{A}(x) = \mathcal{A}^* = \{1\}$  for  $x$  near  $x^*$ , for problem (1).

## 3. Newton's method, $\alpha$ - $\beta$

**TLDR:** For a polynomial system  $F : \mathbb{R}^p \rightarrow \mathbb{R}^p$  and  $z_0 \in \mathbb{R}^p$ , if  $\hat{\alpha}(F, z_0)$  is small enough, there exists a nearby zero of  $F$  to which Newton's method

$$z_{i+1} = z_i - D_F(z_i)^{-1} F(z_i)$$

converge quadratically; see [2].

▷ Some useful quantities

$$\hat{\alpha}(F, z) = \beta(F, z) \mu(F, z) \frac{1}{2} \left( \max_{i=1, \dots, p} d_i \right)^{3/2} \|z\|_{\dagger}^{-1}$$

$$\beta(F, z) = \|D_F(z)^{-1} F(z)\|$$

$$\mu(F, z) = \max\{1, \|F\|_p \|D_F(z)^{-1} \Delta_{(d)}(z)\|\}$$

where  $d_i = \deg F_i$ ,  $\|\cdot\|_F$  is a norm on polynomials, and  $\Delta_{(d)}(z) = \text{Diag}(d_i^{1/2} \|z\|_{\dagger}^{d_i-1})$ , and  $\|z\|_{\dagger} = \sqrt{1 + \|z\|_2^2}$ .

▷ **Proposition** [2]: Consider  $F : \mathbb{R}^p \rightarrow \mathbb{R}^p$  polynomial and  $z_0 \in \mathbb{R}^p$ . If

$$\hat{\alpha}(F, z_0) \leq \frac{1}{4} (13 - 3\sqrt{17}) \approx 0.15, \quad (5)$$

then  $z_0$  is an *approximate zero*:

- $D_F(z_i)$  are invertible i.e., the sequence is well-defined,
- there exists  $\bar{z} \in \mathbb{R}^p$  s.t.  $F(\bar{z}) = 0$ ,  $\|z_0 - \bar{z}\| \leq 2\beta(F, z_0)$ ,

$$\|z_i - \bar{z}\| \leq \left(\frac{1}{2}\right)^{2^i-1} \|z_0 - \bar{z}\|. \quad (6)$$

## The hybrid method

▷ Algorithm outline

- Partial solve of the convex relaxation ( $P_r$ ) e.g., with coord. descent on Burer-Monteiro.
- Extract point  $x_k$  and compute its active set  $\mathcal{A}_k \subseteq \mathcal{I}$ .
- Reduce the problem to

$$\min_x f(x) \quad \text{s.t. } g_i(x) = 0, \quad i \in \mathcal{E} \cup \mathcal{A}_k,$$

and apply Newton to its optimality conditions (2), now *smooth* polynomial equations:

$$F_{\mathcal{A}_k}(x, \lambda) = \left( \begin{array}{c} \nabla f(x) - \sum_{i \in \mathcal{E} \cup \mathcal{A}_k} \lambda_i \nabla g_i(x) \\ g_{\mathcal{E} \cup \mathcal{A}_k}(x) \end{array} \right) = 0,$$

if the  $\alpha$ - $\beta$  test (5) is satisfied.

Then, (6) ensures *superfast convergence*: finding  $(\hat{x}, \hat{\lambda})$   $\epsilon$ -away from  $(x^*, \lambda^*)$  takes  $\lceil 5.82 \rceil = 6$  iterations with  $\epsilon = 2.22 \cdot 10^{-16}$  and  $\|z - z^*\| = 10$ .

- Check global optimality with e.g., the det. hierarchy. Further solve ( $P_r$ ) or increase  $r$  if  $\hat{x}$  is not a global minimizer.

## Algorithm 1 Hybrid algorithm

**Require:**  $(f, g)$ ,  $y_r^0$  initial point for ( $P_r$ )

```

1: for  $k = 0, 1, 2, \dots$  do
2:    $y_r^k \leftarrow$  partial solution of ( $P_r$ ) from point  $y_r^{k-1}$ 
3:    $x^k \leftarrow$  extracted point from SDP iterate  $y_r^k$ 
4:    $\mathcal{A}^k \leftarrow \mathcal{A}(x^k)$ 
5:    $\lambda^k \leftarrow \lambda_{\mathcal{A}^k}(x^k)$ 
6:   if  $\hat{\alpha}(F_{\mathcal{A}^k}, (x^k, \lambda^k)) \leq 0.15$  then
7:      $\hat{x}^k, \hat{\lambda}^k \leftarrow$  limit of Newton's method on  $F_{\mathcal{A}^k}$  from
8:     point  $(x^k, \lambda^k)$ 
9:     if  $(\hat{x}^k, \hat{\lambda}^k)$  satisfies a condition for global opt.
10:    then
11:      return  $\hat{x}^k, \hat{\lambda}^k$ 
12:   else
13:     Further solve ( $P_r$ ) or increase  $r$ 

```

**Theorem** [1, Th. 5.1]: Consider (POP) that admits a unique global minimizer  $x^*$ . Then, algorithm 1 generates  $(x_k, \mathcal{A}_k)$  such that

- $(y_k)$  converges to the *global* minimizer  $x^*$ .

If  $x^*$  is qualified (LICQ and second-order growth), then there exists  $\hat{k}$  such that

- the minimizer active set is identified:  $\mathcal{A}_{\hat{k}} = \mathcal{A}^*$ ,
- Newton's method converges *quadratically* (6),
- and it returns the *globally* optimal primal-dual point  $(\hat{x}^{\hat{k}}, \hat{\lambda}^{\hat{k}})$ .

## Illustrations

▷ Problem: optimal power flow industrial instance

▷ Plots

- fig. 4 shows coordinate descent on ( $P_r$ ) then Newton's method on (POP) after 2, 4, 8, ... epochs
- fig. 5 shows coordinate descent on ( $P_r$ ), Newton's method on (POP), and the hybrid method

▷ Observations

- Coordinate descent on its own can be slow – fig. 5.
- Newton's method converges fast eventually
- Newton's method on its own can be slow – fig. 4 (CD 2 CD 4 CD 8), and fig. 5 (Newton)
- The hybrid method, with the  $\alpha$ - $\beta$  test, improves upon coord. descent and Newton's method – fig. 5
- The limit point attained by a first-then-second order algorithm can vary substantially – fig. 4 right; indeed, Newton's method may be attracted to (first-order stationary) points that are not the global minimizer (see  $x_1, x_2$  in fig. 1), and highlights the need for a global optimality check

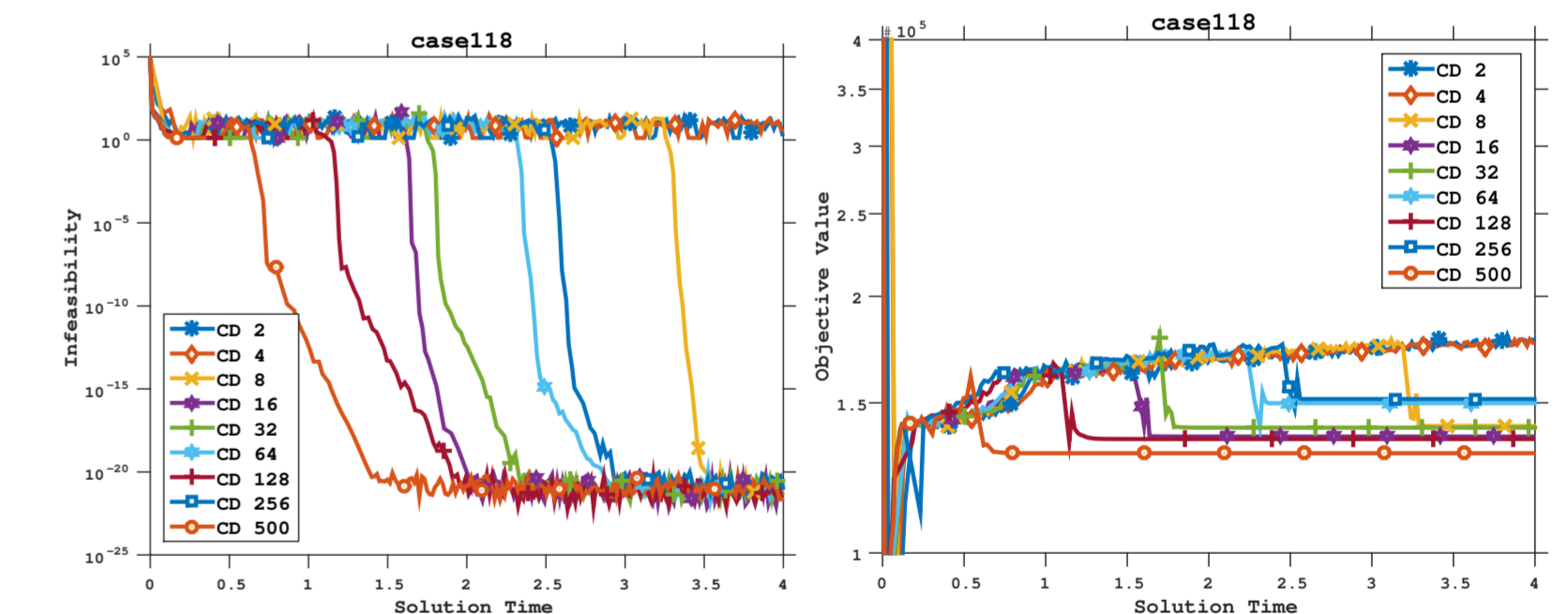


Figure 4. Infeasibility and objective function over time for coordinate descent on ( $P_r$ ) then Newton's method on (POP) after 2, 4, 8, ... epochs.

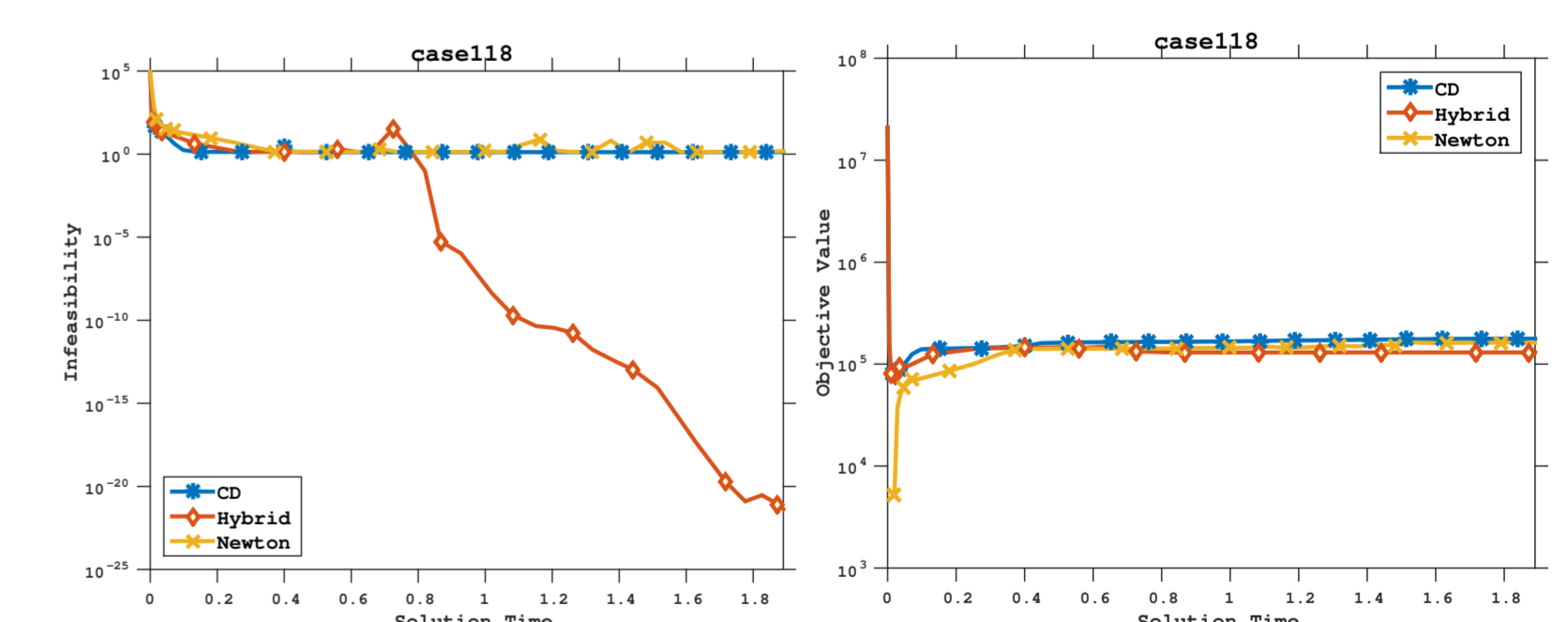


Figure 5. infeasibility and objective function over time for coordinate descent on ( $P_r$ ), Newton's method on (POP), and the hybrid method.



More details in:

[arxiv.org/abs/2305.16122](https://arxiv.org/abs/2305.16122)

## References

- Johannes Aspman, Gilles Bareilles, Vyacheslav Kungurtsev, Jakub Mareček, and Martin Takáč. Hybrid methods in polynomial optimisation, 2023.
- Felipe Cucker and Steve Smale. Complexity estimates depending on condition and round-off error. *Journal of the ACM*, 46(1):113–184, January 1999.
- Jean Bernard Lasserre. *An Introduction to Polynomial and Semi-Algebraic Optimization*. Cambridge University Press, first edition, February 2015.
- Christina Oberlin and Stephen J. Wright. Active Set Identification in Nonlinear Programming. *SIAM Journal on Optimization*, 17(2):577–605, January 2006.